Time-Varying Linear Systems: Identification and Transmission through Unidentified Channels

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Outline

- 1. Introduction: Linear Time-Varying Systems
- 2. Identification under Side Constraints
- 3. Signal Transmission over Unidentified Channels
- 4. Example: Message Transmission with unknown Channel Support

Time-Varying Linear Systems

Time-varying Linear SISO Systems

▷ Single-Input Single-Output Systems (SISO):

Linear time-varying SISO channels are described by operators of $H: L^2(\mathbb{R}) \to L^2(\mathbb{R})$ the form

$$(\mathrm{H}f)(t) = \iint_{\mathbb{R}\times\mathbb{R}} \eta_{\mathrm{H}}(\tau,\nu) \cdot f(t-\tau) \mathrm{e}^{\mathrm{i}2\pi\nu t} \mathrm{d}\nu \mathrm{d}\tau = \iint_{\mathbb{R}\times\mathbb{R}} \eta_{\mathrm{H}}(\tau,\nu) (\mathrm{M}_{\nu}\mathrm{T}_{\tau}f)(t) \mathrm{d}\nu \mathrm{d}\tau$$

with

- Spreading function: $\eta_{H}\mathbb{R}\times\mathbb{R}\rightarrow\mathbb{C}$
- Translation (time-shift) operator: $(T_{\tau}f)(t) = f(t-\tau)$
- Modulation (frequency shift): $M_v(f)(t) = f(t)e^{i2\pi vt}$

Time-varying Linear MIMO Systems

▷ Multiple-Input Multiple-Output Systems (MIMO):

Channels with *N*-inputs and *M*-outputs are characterized by operators $\mathbf{H} : (L^2(\mathbb{R}))^N \to (L^2(\mathbb{R}))^M$ of the form

$$\mathbf{H}\begin{pmatrix}f_{1}\\\vdots\\f_{N}\end{pmatrix}=\begin{pmatrix}\mathbf{H}_{1,1}&\cdots&\mathbf{H}_{1,N}\\\vdots&\vdots\\\mathbf{H}_{M,1}&\cdots&\mathbf{H}_{M,N}\end{pmatrix}\begin{pmatrix}f_{1}\\\vdots\\f_{N}\end{pmatrix}=\begin{pmatrix}\sum_{n=1}^{N}\mathbf{H}_{1,n}f_{n}\\\vdots\\\sum_{n=1}^{N}\mathbf{H}_{M,n}f_{n}\end{pmatrix}$$

Each subchannel $H_{m,n}$ is a TVL SISO system.

Finite-dimensional TVL Channels

▷ The identification problem of TVL systems $\mathbf{H} : (L^2(\mathbb{R}))^N \to (L^2(\mathbb{R}))^M$ can be reduced to a finite-dimensional problem $\mathbf{H} : (\mathbb{C}^L)^N \to (\mathbb{C}^L)^M$

$$\mathbf{H}\begin{pmatrix}\mathbf{x}_{1}\\\vdots\\\mathbf{x}_{N}\end{pmatrix} = \begin{pmatrix}\mathbf{H}_{1,1} & \cdots & \mathbf{H}_{1,N}\\\vdots\\\mathbf{H}_{M,1} & \cdots & \mathbf{H}_{M,N}\end{pmatrix}\begin{pmatrix}\mathbf{x}_{1}\\\vdots\\\mathbf{x}_{N}\end{pmatrix} = \begin{pmatrix}\sum_{n=1}^{N}\mathbf{H}_{1,n}\mathbf{x}_{n}\\\vdots\\\sum_{n=1}^{N}\mathbf{H}_{M,n}\mathbf{x}_{n}\end{pmatrix}$$

wherein each sub-system $\mathbf{H}_{n,m} : \mathbb{C}^L \to \mathbb{C}^L$ has the form

$$\mathbf{H}_{m,n}\mathbf{x} = \sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta_{m,n}(k,\ell) \, \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{x} = \mathbf{G}(\mathbf{x}) \boldsymbol{\eta}$$

with spreading coefficients $\{\eta_{m,n}(k,\ell)\}_{k,\ell=0}^{L-1}$, and with the translation operator $\mathbf{T}: \mathbb{C}^L \to \mathbb{C}^L$ and the modulation operator $\mathbf{M}: \mathbb{C}^L \to \mathbb{C}^L$ given by

$$\mathbf{T}:\begin{pmatrix}x_{0}\\x_{1}\\\vdots\\x_{L-1}\end{pmatrix}\mapsto\begin{pmatrix}x_{L-1}\\x_{0}\\\vdots\\x_{L-2}\end{pmatrix}\quad\text{and}\quad\mathbf{M}:\begin{pmatrix}x_{0}\\x_{1}\\\vdots\\x_{L-1}\end{pmatrix}\mapsto\begin{pmatrix}x_{0}\\x_{1}e^{i\frac{2\pi}{L}\cdot 1}\\\vdots\\x_{L-1}e^{i\frac{2\pi}{L}\cdot(L-1)}\end{pmatrix}$$

Operator Paley–Wiener spaces

 \triangleright Set of all linear operators $\mathbb{C}^L \to \mathbb{C}^L$ is spanned by time-frequency shifts $\mathbf{M}^{\ell} \mathbf{T}^k$

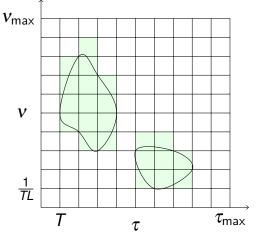
$$\mathscr{L}(\mathbb{C}^{L}) = \left\{ \mathbf{H} = \sum_{k=0}^{L-1} \sum_{\ell=0}^{L-1} \eta(k,\ell) \, \mathbf{M}^{\ell} \mathbf{T}^{k} : \eta(k,\ell) \in \mathbb{C} \text{ for all } (k,\ell) \in \mathbb{Z}_{L} \times \mathbb{Z}_{L} \right\}$$

 \triangleright SISO Operator Paley–Wiener space: For $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$

$$OPW(\Lambda) = \operatorname{span}\left\{\mathbf{M}^{\ell}\mathbf{T}^{k} : (k,\ell) \in \Lambda\right\}$$

▷ MIMO Operator Paley–Wiener space: For $\mathbf{\Lambda} = \{\Lambda_{m,n}\}_{m,n=1}^{M,N}$ with $\Lambda_{m,n} \in \mathbb{Z}_L \times \mathbb{Z}_L$

$$OPW(\mathbf{\Lambda}) = \{\mathbf{H} : \mathbf{H}_{m,n} \in OPW(\Lambda_{m,n})\}$$



Identification under linear side constraints

Identification – SISO

Definition (Identifiable)

The space $OPW(\Lambda)$ is identifiable if and only if there exists an identifier $\mathbf{c} \in \mathbb{C}^L$ such that for each $\mathbf{H} \in OPW(\Lambda)$ the equation

$$\mathbf{y} = \mathbf{H}\mathbf{c} = \sum_{(k,\ell) \in \Lambda} \eta(k,\ell) \, \mathbf{M}^\ell \mathbf{T}^k \, \mathbf{c} = \mathbf{G}(\mathbf{c}) oldsymbol{\eta}$$

is uniquely solvable for $\boldsymbol{\eta} \in \mathbb{C}^{\Lambda}$.

Remark

G(**c**) is Gabor matrix of size $L \times L^2$

$$\mathbf{G}(\mathbf{c}) = \begin{bmatrix} \mathbf{M}^{0} \mathbf{T}^{0} \, \mathbf{c} \ , \mathbf{M}^{0} \mathbf{T}^{1} \, \mathbf{c} \ , \cdots \ , \mathbf{M}^{L-1} \mathbf{T}^{L-1} \, \mathbf{c} \end{bmatrix}$$

Theorem (Identificaltion of SISO Channels) The space $OPW(\Lambda)$ is identifiable if and only if $|\Lambda| \leq L$.

Identification – MIMO

Definition (Identifiable MIMO)

The space $OPW(\Lambda)$ is identifiable if and only if there exist vectors $\mathbf{c} = (\mathbf{c}_1, \mathbf{c}_2, \dots, \mathbf{c}_N) \in (\mathbb{C}^L)^N$ such that for each $\mathbf{H} \in OPW(\Lambda)$ the map $\mathbf{H} \mapsto \mathbf{y} = \mathbf{H}\mathbf{c}$ is injective.

Theorem (Identification of MIMO Channels)

The space $OPW(\Lambda)$ is identifiable if and only if

$$\sum_{n=1}^{N} |\Lambda_{m,n}| \le L \quad \text{for every } m = 1, \dots, M$$

Assumptions

- Subchannels $\mathbf{H}_{m,n}$ and their time-frequency components $\eta_{m,n}(k,\ell)$ are independent
- Information about one channel does not help to identify another.

Linear Constrains Between Spreading Coefficients

▷ Channels **H** in $OPW(\Lambda) \subset \mathscr{L}(\mathbb{C}^L)$:

$$\mathbf{y} = \mathbf{H}\mathbf{c} = \sum_{(k,\ell)\in\Lambda} \eta(k,\ell) \, \mathbf{M}^\ell \mathbf{T}^k \, \mathbf{c} = \mathbf{G}(\mathbf{c}) oldsymbol{\eta}$$

 \triangleright *OPW*(Λ) is identifiable if $|\Lambda| \leq L$.

> Assume linear relations between the spreading coefficients are known

$$\sum_{k,\ell}lpha_{k,\ell}\,\eta(k,\ell)=eta$$
 for some $lpha_{k,\ell},eta\in\mathbb{C}$

Intuition/Question

- Let $A\eta = b$ be a given set of $M \ge 1$ linear independent side constraints.
- Let $OPW_{\mathbf{A},\mathbf{b}}(\Lambda)$ be the set of all $\mathbf{H} \in OPW(\Lambda)$ which satisfy these side constraints.
- $OPW_{\mathbf{A},\mathbf{b}}(\Lambda)$ is identifiable if and only if $|\Lambda| \leq L + M$?

Linear Relations between Time-Frequency Components

> In the SISO case, linear relations between the spreading coefficients of the channel are expressed by

 $\mathbf{b} = \mathbf{A} \boldsymbol{\eta}$

▷ Including the equation for channel identification yields

$$\left[egin{array}{c} {\sf y} \ {\sf b} \end{array}
ight] = \left[egin{array}{c} {\sf G}({\sf c}) \ {\sf A} \end{array}
ight] {m \eta}$$

 \rightarrow More equations for the same number of unknowns should help to identify the channel.

Theorem (One linear side constraint)

Let $\mathbf{H} \in OPW(\Lambda)$ with $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ with $|\Lambda| = L + 1$ and let $\mathbf{a} \in \mathbb{C}^{L+1}$, $\mathbf{a} \neq 0$. There there exists a $\mathbf{c} \in \mathbb{C}^L$ so that

$$\left[\begin{array}{c} \mathbf{G}(\mathbf{c})|_{\Lambda} \\ \mathbf{a}^{*} \end{array} \right]$$

is invertible. Thus the channel coefficients η are identifiable.

Moreover, the set of all such identifiers **c** constitute a dense open subset of \mathbb{C}^{L} .

More Side Constrains

Lemma (No general solution for more than one side constraint)

Let $\mathbf{H} \in OPW(\Lambda)$ with $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ with $|\Lambda| > L + 1$. There exist matrices \mathbf{A} of size $(|\Lambda| - L) \times |\Lambda|$ with $\mathbf{A} \neq 0$ such that there is no $\mathbf{c} \in \mathbb{C}^L$ such that the matrix

$$\left[egin{array}{c} {\bf G}({f c})|_{\Lambda} \ {f A} \end{array}
ight]$$

has full rank.

Theorem (Sufficient condition for identifiably)

Let $\mathbf{H} \in OPW(\Lambda)$ with $\Lambda \in \mathbb{Z}_L \times \mathbb{Z}_L$ of size $|\Lambda| = R > L$. Assume that there exists a subset $\widetilde{\Lambda} \subset \Lambda$ of size L so that (i) $\tau_j(\Lambda) = \tau_j(\widetilde{\Lambda})$ whenever $\tau_j(\Lambda) \neq 0$.

(ii) $\operatorname{ind}(\tau') \neq \operatorname{ind}(\tau(\Lambda))$ for every L-tubel $\tau' \leq \tau(\widetilde{\Lambda})$ of size L different from $\tau(\Lambda)$.

Given any full spark matrix **A** of size $(R - L) \times R$, then there exists an identifier $\mathbf{c} \in \mathbb{C}^{L}$ so that the $R \times R$ matrix

 $\mathbf{G}(\mathbf{c})|_{\Lambda}$

is invertible. Moreover, the set of all such identifiers **c** constitute a dense open subset of
$$\mathbb{C}^L$$
.
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Signal Transmission over Unidentified Channels

Motivation

- ▷ Two step procedure for data transmission over frequency-selective channels
 - 1. Estimate (identify) the channel $\mathbf{H} \in \mathscr{H} \subset \mathscr{L}(\mathbb{C}^{L})$.
 - 2. Transmit data **x** from a certain data set $\mathscr{X} \subset \mathbb{C}^{L}$. Data recovery at the receiver using estimated channel.
- ▷ In time-varying channels, this procedure has to repeated regularly, to update channel state information.
- ▷ In rapidly changing channels, two step procedure becomes more and more inefficient.

Transmission through unidentified channel

- ▷ Combine channel identification and signal recovery.
- \triangleright Transmission scheme $\mathbf{y} = \mathbf{H}(\mathbf{x} + \mathbf{c})$ with
 - $\mathbf{H} \in \mathscr{H} \subset \mathscr{L}(\mathbb{C}^{L})$ is a unknown channel from a known subset \mathscr{H} .
 - $-\mathbf{x} \in \mathscr{X}$ is the data signal from a certain data set $\mathscr{X} \subset \mathbb{C}^{L}$.
 - $-\mathbf{c} \in \mathbb{C}^{L}$ is a pilot signal (designed based on the knowledge of \mathscr{H} and \mathscr{X}).

Problem

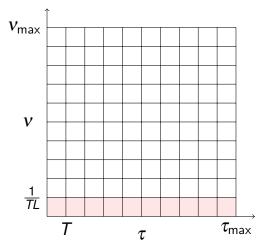
Find (necessary and/or sufficient) conditions on \mathscr{H} and \mathscr{X} such that there exists a pilot $\mathbf{c} \in \mathbb{C}^{L}$ such that every $\mathbf{x} \in \mathscr{X}$ can uniquely be recovered form $\mathbf{y} = \mathbf{H}(\mathbf{x} + \mathbf{c})$ for any unknown channel $\mathbf{H} \in \mathscr{H}$.

Relation to Blind Deconvolution

$$\triangleright \ \mathscr{H} = OPW(\Lambda) \text{ with } \Lambda\{(0,0),(1,0),\ldots,(L,0)\} \subset \mathbb{Z}_L \times \mathbb{Z}_L$$

$$\mathbf{y} = \mathbf{H}\mathbf{c} = \sum_{k=0}^{L-1} \eta(k,0) \mathbf{T}^k \mathbf{x}$$

▷ Recover **x** ∈ X ∈ C^L and **η** from **y** (without knowing η).
▷ **x** and **η** assumed to be sparse.



Conditions for Data Recovery and/or Channel Identification

Natural Conditions

(I) Identifiability of \mathscr{H} : The map $\mathbf{H} \mapsto \mathbf{Hc}$ is injective on \mathscr{H} .

(R) Recovery condition for known channel: Every $\mathbf{H} \in \mathscr{H}$ is injective on \mathscr{X} .

Subsets of \mathbb{C}^{L}

 $\mathscr{H}\mathbf{C} = \{\mathbf{H}\mathbf{C} : \mathbf{H} \in \mathscr{H}\}$: All possible output vectors for the pilot **c**.

 $\mathscr{H}\mathscr{X} = \{\mathbf{H}\mathbf{x} : \mathbf{H} \in \mathscr{H}, \mathbf{x} \in \mathscr{X}\}$: Possible output of arbitrary data vector in \mathscr{X} .

Further Conditions

- (i) $\operatorname{span}{\mathscr{H}\mathbf{C}} \cap \operatorname{span}{\mathscr{H}\mathscr{X}} = \{\mathbf{0}\}$: Isolate **Hc** and **Hx** from the channel output $\mathbf{y} = \mathbf{H}(\mathbf{c} + \mathbf{x})$.
- (ii)
- (iii) $\mathbf{H}(\mathbf{x} + \mathbf{c}) = \mathbf{H}'(\mathbf{x}' + \mathbf{c})$ implies $\mathbf{x} = \mathbf{x}'$
- $H(\mathscr{X} + c) \cap H'(\mathscr{X} + c)$ for every $H \neq H'$ in \mathscr{H} : Identify H from y = H(c + x) with unknown $x \in \mathscr{X}$.
 - : Guarantees exact recovery of $\mathbf{x} \in \mathscr{X}$ but not identification **H**

$$(i) \stackrel{(I)}{\Longrightarrow} (ii) \stackrel{(R)}{\Longrightarrow} (iii)$$

(1)

Degrees of Freedom

 $\triangleright \mathscr{H} = OPW(\Lambda) \subset \mathscr{L}(\mathbb{C}^L)$ is a linear subspace of dimension $|\Lambda|$.

- ▷ Assume $\mathscr{X} \subset \mathbb{C}^L$ is a linear subspace of dimension *K*.
- > Counting degrees of freedom, we must have

$$|\Lambda| + K \leq L$$
 and $|\Lambda| \leq L$.

as a necessary condition for exact recovery of $\mathbf{H} \in \mathscr{H}$ and $\mathbf{x} \in \mathscr{X}$.

- ▷ Without identifying **H**, we may get $K > L |\Lambda| \implies$ how?
- \triangleright When we get equality in (1)?

Example: 1-Dimensional Signal Space

 \triangleright For a given (and known) $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$, we consider the operator Paley–Wiener space

$$\mathscr{H} = OPW(\Lambda) = \Big\{ \mathsf{H} = \sum_{k,\ell \in \Lambda} \eta(k,\ell) \, \mathsf{M}^\ell \mathsf{T}^k \; : \; \eta(k,\ell) \in \mathbb{C} \Big\}. \subset \mathscr{L}(\mathbb{C}^L)$$

▷ Assume a 1-dimensional signal space $\mathscr{X} = \operatorname{span}\{\mathbf{v}\} \subset \mathbb{C}^L$ for some $\mathbf{v} \in \mathbb{C}^L$.

▷ For $x = u \mathbf{v} \in \mathscr{X}$, with $u \in \mathbb{C}$, the received signal is

$$\mathbf{y} = \mathbf{H}(\mathbf{x} + \mathbf{c}) = u \mathbf{H}(\mathbf{v}) + \mathbf{H}(\mathbf{c}) = \mathbf{G}(\mathbf{c})|_{\wedge} \boldsymbol{\eta} + u \mathbf{G}(\mathbf{v})|_{\wedge} \boldsymbol{\eta}$$

 $\triangleright \text{ Separation of } \mathbf{Hx} \text{ and } \mathbf{Hc} \text{ from } \mathbf{y} \Longrightarrow \operatorname{span}\{\mathscr{H}\mathbf{c}\} \cap \operatorname{span}\{\mathscr{H}\mathbf{v}\} = \{\mathbf{0}\} \Longrightarrow \operatorname{rang}[\mathbf{G}(\mathbf{c})|_{\Lambda}] \perp \operatorname{rang}[\mathbf{G}(\mathbf{v})|_{\Lambda}].$

For rang $[\mathbf{G}(\mathbf{v})|_{\Lambda}] = \operatorname{span}\{\mathbf{a}\}\$ is a 1-dimensional subspace, we have the following result.

Theorem Let $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ with $1 \leq |\Lambda| \leq L-1$ and let $\mathbf{a} \in \mathbb{C}^L \setminus \{0\}$ be arbitrary. There exists a vector $\mathbf{c} \in \mathbb{C}^L$ such that the $L \times |\Lambda| + 1$ matrix $[\mathbf{G}(\mathbf{c})|_{\Lambda}, \mathbf{a}]$ has full rank.

General Construction

▷ The set of all time-frequency shifts $\{\mathbf{M}^{\ell}\mathbf{T}^k\}_{\ell,k=0}^{L-1}$ can be separated into L+1 commutative subgroups

$$\mathcal{G}_{s} = \left\{ \mathbf{M}^{2rk} \mathbf{T}^{k} : k = 0, 1, \dots, L - 1 \right\}, \qquad s = 0, 1, \dots, L - 1$$
$$\mathcal{G}_{L} = \left\{ \mathbf{M}^{k} : k = 0, 1, \dots, L - 1 \right\}$$

 \triangleright Each commutative subgroup \mathscr{G}_s posses a set of common eigenvectors (i.e. the *chirp sequences*)

$$e_{s}(\ell): \ell = 0, 1, \dots, L-1$$

which forms on orthogonal basis for \mathbb{C}^{L} .

Theorem

Let $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ such that there exists an $s \in \{0, 1, ..., L\}$ so that $\Lambda \subset \mathscr{G}_s$. There exists a subspace $\mathscr{X} = \operatorname{span}\{\mathbf{e}_s(1), ..., \mathbf{e}_s(K)\} \subset \mathbb{C}^L$ of dimension $K = L - |\Lambda|$ and a $\mathbf{c} \in \mathbb{C}^L$ such that

$$\operatorname{span}\{\mathscr{H}\mathbf{C}\}\cap\operatorname{span}\{\mathscr{H}\mathscr{X}\}=\{\mathbf{0}\}$$

and such that Conditions (I) and (R) are satisfied.

Maximum Size of Data Subspace

▷ Given a subspace \mathscr{X} of dimension *K*, it is desirable that the dimension of span{ $\mathscr{H}\mathscr{X}$ } is again *K*.

▷ For for a one-dimensional $\mathscr{X} = \operatorname{span}\{\mathbf{v}\}$, we have $\dim(\operatorname{span}\{\mathscr{H}\mathscr{X}\}) = \operatorname{rank}(\mathbf{G}(\mathbf{v})|_{\Lambda})$.

Question: Given a support set $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ of $\mathscr{H} = OPW(\Lambda)$ with $|\Lambda| \leq L$. What is the minimum rank of $\mathbf{G}(\mathbf{v})|_{\Lambda}$ for \mathbf{v} varying in \mathbb{C}^L ?

Theorem

Let $L \ge 3$ be an odd integer and $\Lambda \subset \mathbb{Z}_L \times \mathbb{Z}_L$ with $|\Lambda| \le L$. Then

 $\min_{\boldsymbol{v}\in\mathbb{C}^{L}\setminus\{\boldsymbol{0}\}}\mathrm{rank}\left(\boldsymbol{\mathsf{G}}(\boldsymbol{v})|_{\Lambda}\right)\leq \textit{N}(\Lambda)$

with

$$N(\Lambda) = 1 + \min_{s \in \{0,1,\dots,L\}} \min \left\{ |I| : I \subset \mathbb{Z}_L \times \mathbb{Z}_L \text{ with } \Lambda \subset I + \mathscr{G}_s \right\},$$

and where $I + \mathscr{G}_s = \{x + y : x \in I, y \in \mathscr{G}_s\}.$

Example with Unknown Channel Support

Channel Model

 \triangleright

▷ Time continuous Time-Varying-Linear SISO Channel $H: L^2(\mathbb{R}) \to L^2(\mathbb{R})$:

$$g(t) = (\mathrm{H}f)(t) = \iint_{\mathbb{R}\times\mathbb{R}} \eta_{\mathrm{H}}(\tau, \nu) \cdot f(t-\tau) \mathrm{e}^{\mathrm{i}2\pi\nu t} \mathrm{d}\nu \mathrm{d}\tau = \iint_{\mathbb{R}\times\mathbb{R}} \eta_{\mathrm{H}}(\tau, \nu) (\mathrm{M}_{\nu}\mathrm{T}_{\tau}f)(t) \mathrm{d}\nu \mathrm{d}\tau$$

▷ Time discrete Time-Varying-Linear SISO Channel $H: L^2(\mathbb{R}) \to L^2(\mathbb{R})$:

$$g(t) = (\mathrm{H}f)(t) = \sum_{k=0}^{K-1} \sum_{\ell=0}^{M-1} \eta_{\mathrm{H}}(k,\ell) \cdot f(t-k\Delta\tau) e^{\mathrm{i}2\pi\ell\Delta\nu t} = \sum_{k=0}^{K-1} \sum_{\ell=0}^{M-1} \eta_{\mathrm{H}}(k,\ell) \cdot f(t-kT) e^{\mathrm{i}\frac{2\pi}{TL}\ell t}$$
Rectification of the channel support region:
with $\Delta\tau = T$ and $\Delta\nu = \frac{1}{TL}$ for some prime $L \ge 5$.

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Transmitted Signal

▷ Transmitted signal: Delta train followed by a guard interval

$$f(t) = \sum_{m=0}^{2L-1} x_m \delta(t - mT) \quad \text{with} \quad x_m = \begin{cases} \text{data symbol} : & 0 \le m \le L-1 \\ 0 & : & L \le m \le 2L-1 \end{cases}$$

- ▷ Received signal:
 - 1. Sampling at rate 1/T: $g_n = (Hf)(nT), n = 0, 1, ..., 2L 1.$
 - 2. Add two consecutive blocks of length *L*: $y_n = g_n + g_{n+L}$, n = 0, 1, ..., L 1.
- ▷ Write in vector form:

$$\mathbf{y} = \mathbf{H}\mathbf{x} = \sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta(k, l) \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{x}$$

with $\mathbf{x} = (x_0, \dots, x_{L-1})^T$ and $\mathbf{y} = (x_0, \dots, x_{L-1})^T$ and with $\mathbf{H} \in OPW(\Lambda)$.

Remark: Using a periodic weighted delta train

$$f(t) = \sum_{p \in \mathbb{Z}} \sum_{m=0}^{2L-1} x_m \delta(t - mT - pLT)$$

results in a similar expression as in (2) but requires periodic data. Volker Pohl (TUM) | Time-Varying Linear Systems | CoSIP 2020 (2)

Transmission Scheme

- ▷ Aim: We want to transmit a message $\gamma = \{\gamma_1, ..., \gamma_Q\} \subset \mathbb{C}$ of size $Q \leq L$ over the channel $\mathbf{H} \in OPW(\Lambda)$ and recover γ at the receiver without knowing \mathbf{H} and without knowing the channel support Λ in advance.
- > Data symbols: The data symbols a sum of a pilot and the actual message

$$\mathbf{x} = \mathbf{c} + \sum_{q=1}^Q \gamma_q \; \mathbf{e}_q$$

with the pilot signal **c** which is chosen to be an *Alltop sequence* and \mathbf{e}_q are particular *chirp sequences*

$$\mathbf{c}(n) = \frac{1}{\sqrt{L}} \exp\left(\mathrm{i}\frac{2\pi}{L}n^3\right) \quad \text{and} \quad \mathbf{e}_{mL+r}(n) = \frac{1}{\sqrt{L}} \exp\left(\mathrm{i}\frac{2\pi}{L}\left[r+mn+rn^2\right]\right), \qquad n \in \mathbb{Z}_L$$

▷ Received signal:

$$\mathbf{y} = \mathbf{H}\mathbf{x} = \sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta(k,l) \, \mathbf{M}^\ell \mathbf{T}^k \mathbf{x} = \mathbf{G}(\mathbf{x}) \boldsymbol{\eta} \; = \; \mathbf{G}(\mathbf{c}) \boldsymbol{\eta} + \mathbf{U}\mathbf{s} = \mathbf{\Phi} \left(egin{array}{c} m{\eta} \\ \mathbf{s} \end{array}
ight)$$

with a $L \times 2L^2$ measurement matrix $\mathbf{\Phi} = [\mathbf{G}(\mathbf{c}), \mathbf{U}]$ and a sparse vector $\mathbf{s} = f(\boldsymbol{\eta}, \boldsymbol{\gamma})$. \triangleright Recovery of the message: Solve CS problem (3), then determine $\boldsymbol{\gamma} = g(\boldsymbol{\eta}, \mathbf{s})$. (3)

Compressive Sampling Problem

 \triangleright Compressive Sampling Problem of size $L \times 2L^2$

$$\mathbf{y} = \mathbf{\Phi} \left(egin{array}{c} m{\eta} \\ \mathbf{s} \end{array}
ight)$$
 with $\left| \operatorname{supp} (m{\eta}, \mathbf{s})^{\mathrm{T}} \right| \leq (1 + Q) |\Lambda|$

Lemma

The coherence of the measurement matrix $\Phi = [\mathbf{G}(\mathbf{c}), \mathbf{U}]$ is upper bounded by $\mu(\Phi) \leq \frac{2}{\sqrt{L}}$.

Remark: Welsh bound is $\frac{1}{\sqrt{L+1}}$.

Theorem

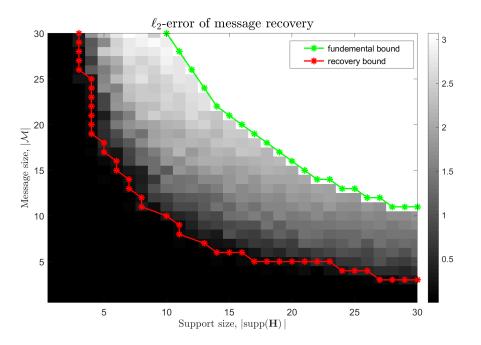
Let $\mathbf{H} \in OPW(\Lambda) \subset \mathscr{L}(\mathbb{C}^L)$ be an unknown channel with unknown support set Λ where $L \ge 5$ is a prime number. For

$$Q \leq rac{\sqrt{L}}{4|\Lambda|} - \frac{1}{2}$$

any message $\gamma \in \mathbb{C}^Q$ can be transmitted over **H** and recovered at the receiver.

Numerical Experiment

- *L* = 307
- OMP
- average over 100 channels (random support, Gaussian coefficients)
- red line: 1% rel. error
- simulation much better than bound



Summary

Current and Future Work

▷ Identification of SISO and MIMO TVL Channels:

- Identification of stochastic channels and stochastic sequences
- Conditions on support of the covariance of the spreading function η .
- Linear side constraints in terms of covariance.

▷ Transmission over Unidentified Channels:

- Stochastic encoding, RIP
- Scheme which maximum transmission rate
- Continuous time setting

Related Publications

- D. G. Lee, G. E. Pfander, V. Pohl, W. Zhou "Identification of Channels with Single and Multiple Inputs and Outputs under Linear Constraints," *Linear Algebra Appl.*, vol. 581 (Nov. 2019), 435 470.
- D. G. Lee, G. E. Pfander, V. Pohl "Signal transmission through an unidentified channel," *13th Intern. Conf. on Sampling Theory and Applications (SampTA)*, Bordeaux, France, July 2019.
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