# Time-Varying Linear Systems: Identification and Transmission through Unidentified Channels 

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## Outline

1. Introduction: Linear Time-Varying Systems
2. Identification under Side Constraints
3. Signal Transmission over Unidentified Channels
4. Example: Message Transmission with unknown Channel Support

## Time-Varying Linear Systems

## Time-varying Linear SISO Systems

$\triangleright$ Single-Input Single-Output Systems (SISO):
Linear time-varying SISO channels are described by operators of $H: L^{2}(\mathbb{R}) \rightarrow L^{2}(\mathbb{R})$ the form

$$
(\mathrm{H} f)(t)=\iint_{\mathbb{R} \times \mathbb{R}} \eta_{\mathrm{H}}(\tau, v) \cdot f(t-\tau) \mathrm{e}^{\mathrm{i} 2 \pi v t} \mathrm{~d} v \mathrm{~d} \tau=\iint_{\mathbb{R} \times \mathbb{R}} \eta_{\mathrm{H}}(\tau, v)\left(\mathrm{M}_{v} \mathrm{~T}_{\tau} f\right)(t) \mathrm{d} v \mathrm{~d} \tau
$$

with

- Spreading function: $\eta_{\mathrm{H}} \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{C}$
- Translation (time-shift) operator: $\left(\mathrm{T}_{\tau} f\right)(t)=f(t-\tau)$
- Modulation (frequency shift): $\mathrm{M}_{v}(f)(t)=f(t) \mathrm{e}^{\mathrm{i} 2 \pi v t}$


## Time-varying Linear MIMO Systems

$\triangleright$ Multiple-Input Multiple-Output Systems (MIMO):
Channels with $N$-inputs and $M$-outputs are characterized by operators $\mathbf{H}:\left(L^{2}(\mathbb{R})\right)^{N} \rightarrow\left(L^{2}(\mathbb{R})\right)^{M}$ of the form

$$
\mathbf{H}\left(\begin{array}{c}
f_{1} \\
\vdots \\
f_{N}
\end{array}\right)=\left(\begin{array}{ccc}
\mathrm{H}_{1,1} & \cdots & \mathrm{H}_{1, N} \\
\vdots & & \vdots \\
\mathrm{H}_{M, 1} & \cdots & \mathrm{H}_{M, N}
\end{array}\right)\left(\begin{array}{c}
f_{1} \\
\vdots \\
f_{N}
\end{array}\right)=\left(\begin{array}{c}
\sum_{n=1}^{N} \mathrm{H}_{1, n} f_{n} \\
\vdots \\
\sum_{n=1}^{N} \mathrm{H}_{M, n} f_{n}
\end{array}\right) .
$$

Each subchannel $H_{m, n}$ is a TVL SISO system.

## Finite-dimensional TVL Channels

$\triangleright$ The identification problem of TVL systems $\mathbf{H}:\left(L^{2}(\mathbb{R})\right)^{N} \rightarrow\left(L^{2}(\mathbb{R})\right)^{M}$ can be reduced to a finite-dimensional problem $\mathbf{H}:\left(\mathbb{C}^{L}\right)^{N} \rightarrow\left(\mathbb{C}^{L}\right)^{M}$

$$
\mathbf{H}\left(\begin{array}{c}
\mathbf{x}_{1} \\
\vdots \\
\mathbf{x}_{N}
\end{array}\right)=\left(\begin{array}{ccc}
\mathbf{H}_{1,1} & \cdots & \mathbf{H}_{1, N} \\
\vdots & & \vdots \\
\mathbf{H}_{M, 1} & \cdots & \mathbf{H}_{M, N}
\end{array}\right)\left(\begin{array}{c}
\mathbf{x}_{1} \\
\vdots \\
\mathbf{x}_{N}
\end{array}\right)=\left(\begin{array}{c}
\sum_{n=1}^{N} \mathbf{H}_{1, n} \mathbf{x}_{n} \\
\vdots \\
\sum_{n=1}^{N} \mathbf{H}_{M, n} \mathbf{x}_{n}
\end{array}\right) .
$$

wherein each sub-system $\mathbf{H}_{n, m}: \mathbb{C}^{L} \rightarrow \mathbb{C}^{L}$ has the form

$$
\mathbf{H}_{m, n} \mathbf{x}=\sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta_{m, n}(k, \ell) \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{x}=\mathbf{G}(\mathbf{x}) \boldsymbol{\eta}
$$

with spreading coefficients $\left\{\eta_{m, n}(k, \ell)\right\}_{k, \ell=0}^{L-1}$, and with the translation operator $\mathbf{T}: \mathbb{C}^{L} \rightarrow \mathbb{C}^{L}$ and the modulation operator $\mathbf{M}: \mathbb{C}^{L} \rightarrow \mathbb{C}^{L}$ given by

$$
\mathbf{T}:\left(\begin{array}{c}
x_{0} \\
x_{1} \\
\vdots \\
x_{L-1}
\end{array}\right) \mapsto\left(\begin{array}{c}
x_{L-1} \\
x_{0} \\
\vdots \\
x_{L-2}
\end{array}\right) \quad \text { and } \quad \mathbf{M}:\left(\begin{array}{c}
x_{0} \\
x_{1} \\
\vdots \\
x_{L-1}
\end{array}\right) \mapsto\left(\begin{array}{c}
x_{0} \\
x_{1} \mathrm{e}^{\mathrm{i} \frac{2 \pi}{L} \cdot 1} \\
\vdots \\
x_{L-1} \mathrm{e}^{\mathrm{i} \frac{2 \pi}{L} \cdot(L-1)}
\end{array}\right)
$$

## Operator Paley-Wiener spaces

$\triangleright$ Set of all linear operators $\mathbb{C}^{L} \rightarrow \mathbb{C}^{L}$ is spanned by time-frequency shifts $\mathbf{M}^{\ell} \mathbf{T}^{k}$

$$
\mathscr{L}\left(\mathbb{C}^{L}\right)=\left\{\mathbf{H}=\sum_{k=0}^{L-1} \sum_{\ell=0}^{L-1} \eta(k, \ell) \mathbf{M}^{\ell} \mathbf{T}^{k}: \eta(k, \ell) \in \mathbb{C} \text { for all }(k, \ell) \in \mathbb{Z}_{L} \times \mathbb{Z}_{L}\right\}
$$

$\triangleright$ SISO Operator Paley-Wiener space: For $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$

$$
O P W(\Lambda)=\operatorname{span}\left\{\mathbf{M}^{\ell} \mathbf{T}^{\kappa}:(k, \ell) \in \Lambda\right\}
$$

$\triangleright$ MIMO Operator Paley-Wiener space: For $\boldsymbol{\Lambda}=\left\{\Lambda_{m, n}\right\}_{m, n=1}^{M, N}$ with $\Lambda_{m, n} \in \mathbb{Z}_{L} \times \mathbb{Z}_{L}$

$$
O P W(\boldsymbol{\Lambda})=\left\{\mathbf{H}: \mathbf{H}_{m, n} \in O P W\left(\Lambda_{m, n}\right)\right\}
$$



## Identification under linear side constraints

## Identification - SISO

## Definition (Identifiable)

The space $\operatorname{OPW}(\Lambda)$ is identifiable if and only if there exists an identifier $\mathbf{c} \in \mathbb{C}^{L}$ such that for each $\mathbf{H} \in O P W(\Lambda)$ the equation

$$
\mathbf{y}=\mathbf{H} \mathbf{c}=\sum_{(k, \ell) \in \Lambda} \eta(k, \ell) \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{c}=\mathbf{G}(\mathbf{c}) \boldsymbol{\eta}
$$

is uniquely solvable for $\boldsymbol{\eta} \in \mathbb{C}^{\wedge}$.
Remark
$\mathbf{G}(\mathbf{c})$ is Gabor matrix of size $L \times L^{2}$

$$
\mathbf{G}(\mathbf{c})=\left[\mathbf{M}^{0} \mathbf{T}^{0} \mathbf{c}, \mathbf{M}^{0} \mathbf{T}^{1} \mathbf{c}, \cdots, \mathbf{M}^{L-1} \mathbf{T}^{L-1} \mathbf{c}\right]
$$

## Theorem (Identificaltion of SISO Channels)

The space $\operatorname{OPW}(\Lambda)$ is identifiable if and only if $|\Lambda| \leq L$.

## Identification - MIMO

## Definition (Identifiable MIMO)

The space $\operatorname{OPW}(\boldsymbol{\Lambda})$ is identifiable if and only if there exist vectors $\mathbf{c}=\left(\mathbf{c}_{1}, \mathbf{c}_{2}, \ldots, \mathbf{c}_{N}\right) \in\left(\mathbb{C}^{L}\right)^{N}$ such that for each $\mathbf{H} \in \operatorname{OPW}(\boldsymbol{\Lambda})$ the $\operatorname{map} \mathbf{H} \mapsto \mathbf{y}=\mathbf{H c}$ is injective.
Theorem (Identification of MIMO Channels)
The space $\operatorname{OPW}(\Lambda)$ is identifiable if and only if

$$
\sum_{n=1}^{N}\left|\Lambda_{m, n}\right| \leq L \quad \text { for every } m=1, \ldots, M
$$

## Assumptions

- Subchannels $\mathbf{H}_{m, n}$ and their time-frequency components $\eta_{m, n}(k, \ell)$ are independent
- Information about one channel does not help to identify another.


## Linear Constrains Between Spreading Coefficients

$\triangleright$ Channels H in $\mathrm{OPW}(\Lambda) \subset \mathscr{L}\left(\mathbb{C}^{L}\right)$ :

$$
\mathbf{y}=\mathbf{H} \mathbf{c}=\sum_{(k, \ell) \in \Lambda} \eta(k, \ell) \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{c}=\mathbf{G}(\mathbf{c}) \boldsymbol{\eta}
$$

$\triangleright \operatorname{OPW}(\Lambda)$ is identifiable if $|\Lambda| \leq L$.
$\triangleright$ Assume linear relations between the spreading coefficients are known

$$
\sum_{k, \ell} \alpha_{k, \ell} \eta(k, \ell)=\beta \quad \text { for some } \quad \alpha_{k, \ell}, \beta \in \mathbb{C}
$$

Intuition/Question

- Let $\mathbf{A} \boldsymbol{\eta}=\mathbf{b}$ be a given set of $M \geq 1$ linear independent side constraints.
- Let $O P W_{\mathbf{A}, \mathbf{b}}(\Lambda)$ be the set of all $\mathbf{H} \in O P W(\Lambda)$ which satisfy these side constraints.
- $O P W_{\mathrm{A}, \mathrm{b}}(\Lambda)$ is identifiable if and only if $|\Lambda| \leq L+M$ ?


## Linear Relations between Time-Frequency Components

$\triangleright$ In the SISO case, linear relations between the spreading coefficients of the channel are expressed by

$$
\mathbf{b}=\mathbf{A} \boldsymbol{\eta}
$$

$\triangleright$ Including the equation for channel identification yields

$$
\left[\begin{array}{l}
\mathbf{y} \\
\mathbf{b}
\end{array}\right]=\left[\begin{array}{c}
\mathbf{G}(\mathbf{c}) \\
\mathbf{A}
\end{array}\right] \boldsymbol{\eta}
$$

$\rightarrow$ More equations for the same number of unknowns should help to identify the channel.
Theorem (One linear side constraint)
Let $\mathbf{H} \in O P W(\Lambda)$ with $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ with $|\Lambda|=L+1$ and let $\mathbf{a} \in \mathbb{C}^{L+1}, \mathbf{a} \neq 0$. There there exists a $\mathbf{c} \in \mathbb{C}^{L}$ so that

$$
\left[\begin{array}{c}
\mathbf{G}(\mathbf{c}) \mid \wedge \\
\mathbf{a}^{*}
\end{array}\right]
$$

is invertible. Thus the channel coefficients $\eta$ are identifiable.
Moreover, the set of all such identifiers constitute a dense open subset of $\mathbb{C}^{L}$.

## More Side Constrains

Lemma (No general solution for more than one side constraint)
Let $\mathbf{H} \in O P W(\Lambda)$ with $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ with $|\Lambda|>L+1$. There exist matrices $\mathbf{A}$ of size $(|\Lambda|-L) \times|\Lambda|$ with $\mathbf{A} \neq 0$ such that there is no $\mathbf{c} \in \mathbb{C}^{L}$ such that the matrix

$$
\left[\begin{array}{c}
\left.\mathbf{G}(\mathbf{c})\right|_{\Lambda} \\
\mathbf{A}
\end{array}\right]
$$

has full rank.

## Theorem (Sufficient condition for identifiably)

Let $\mathbf{H} \in \operatorname{OPW}(\Lambda)$ with $\Lambda \in \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ of size $|\Lambda|=R>L$. Assume that there exists a subset $\widetilde{\Lambda} \subset \Lambda$ of size $L$ so that
(i) $\tau_{j}(\Lambda)=\tau_{j}(\widetilde{\Lambda})$ whenever $\tau_{j}(\Lambda) \neq 0$.
(ii) $\operatorname{ind}\left(\tau^{\prime}\right) \neq \operatorname{ind}(\tau(\Lambda))$ for every L-tubel $\tau^{\prime} \preceq \tau(\widetilde{\Lambda})$ of size $L$ different from $\tau(\Lambda)$.

Given any full spark matrix $\mathbf{A}$ of size $(R-L) \times R$, then there exists an identifier $\mathbf{c} \in \mathbb{C}^{L}$ so that the $R \times R$ matrix

$$
\left[\begin{array}{c}
\left.\mathbf{G}(\mathbf{c})\right|_{\Lambda} \\
\mathbf{A}
\end{array}\right]
$$

is invertible. Moreover, the set of all such identifiers constitute a dense open subset of $\mathbb{C}^{L}$.

# Signal Transmission over Unidentified Channels 

## Motivation

$\triangleright$ Two step procedure for data transmission over frequency-selective channels

1. Estimate (identify) the channel $\mathbf{H} \in \mathscr{H} \subset \mathscr{L}\left(\mathbb{C}^{L}\right)$.
2. Transmit data $\mathbf{x}$ from a certain data set $\mathscr{X} \subset \mathbb{C}^{L}$. Data recovery at the receiver using estimated channel.
$\triangleright$ In time-varying channels, this procedure has to repeated regularly, to update channel state information.
$\triangleright$ In rapidly changing channels, two step procedure becomes more and more inefficient.

## Transmission through unidentified channel

$\triangleright$ Combine channel identification and signal recovery.
$\triangleright$ Transmission scheme $\mathbf{y}=\mathbf{H}(\mathbf{x}+\mathbf{c})$ with
$-\mathbf{H} \in \mathscr{H} \subset \mathscr{L}\left(\mathbb{C}^{L}\right)$ is a unknown channel from a known subset $\mathscr{H}$.
$-\mathbf{x} \in \mathscr{X}$ is the data signal from a certain data set $\mathscr{X} \subset \mathbb{C}^{L}$.
$-\mathbf{c} \in \mathbb{C}^{L}$ is a pilot signal (designed based on the knowledge of $\mathscr{H}$ and $\mathscr{X}$ ).

## Problem

Find (necessary and/or sufficient) conditions on $\mathscr{H}$ and $\mathscr{X}$ such that there exists a pilot $\mathbf{c} \in \mathbb{C}^{L}$ such that every $\mathbf{x} \in \mathscr{X}$ can uniquely be recovered form $\mathbf{y}=\mathbf{H}(\mathbf{x}+\mathbf{c})$ for any unknown channel $\mathbf{H} \in \mathscr{H}$.

## Relation to Blind Deconvolution

$\triangleright \mathscr{H}=\operatorname{OPW}(\Lambda)$ with $\Lambda\{(0,0),(1,0), \ldots,(L, 0)\} \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$

$$
\mathbf{y}=\mathbf{H} \mathbf{c}=\sum_{k=0}^{L-1} \eta(k, 0) \mathbf{T}^{k} \mathbf{x}
$$

$\triangleright$ Recover $\mathbf{x} \in \mathscr{X} \in \mathbb{C}^{L}$ and $\boldsymbol{\eta}$ from $\mathbf{y}$ (without knowing $\eta$ ).
$\triangleright \mathbf{x}$ and $\boldsymbol{\eta}$ assumed to be sparse.


## Conditions for Data Recovery and/or Channel Identification

## Natural Conditions

(I) Identifiability of $\mathscr{H}$ : The map $\mathbf{H} \mapsto \mathbf{H c}$ is injective on $\mathscr{H}$.
(R) Recovery condition for known channel: Every $\mathbf{H} \in \mathscr{H}$ is injective on $\mathscr{X}$.

Subsets of $\mathbb{C}^{L}$

$$
\begin{array}{ll}
\mathscr{H} \mathbf{c}=\{\mathbf{H} \mathbf{c}: \mathbf{H} \in \mathscr{H}\} & : \text { All possible output vectors for the pilot } \mathbf{c} . \\
\mathscr{H} \mathscr{X}=\{\mathbf{H} \mathbf{x}: \mathbf{H} \in \mathscr{H}, \mathbf{x} \in \mathscr{X}\} & : \text { Possible output of arbitrary data vector in } \mathscr{X} .
\end{array}
$$

Further Conditions
(i) $\operatorname{span}\{\mathscr{H} \mathbf{c}\} \cap \operatorname{span}\{\mathscr{H} \mathscr{X}\}=\{0\} \quad$ : Isolate $\mathbf{H c}$ and $\mathbf{H x}$ from the channel output $\mathbf{y}=\mathbf{H}(\mathbf{c}+\mathbf{x})$.
(ii) $\mathbf{H}(\mathscr{X}+\mathbf{c}) \cap \mathbf{H}^{\prime}(\mathscr{X}+\mathbf{c})$ for every $\mathbf{H} \neq \mathbf{H}^{\prime}$ in $\mathscr{H}$ : Identify $\mathbf{H}$ from $\mathbf{y}=\mathbf{H}(\mathbf{c}+\mathbf{x})$ with unknown $\mathbf{x} \in \mathscr{X}$.
(iii) $\mathbf{H}(\mathbf{x}+\mathbf{c})=\mathbf{H}^{\prime}\left(\mathbf{x}^{\prime}+\mathbf{c}\right)$ implies $\mathbf{x}=\mathbf{x}^{\prime} \quad$ : Guarantees exact recovery of $\mathbf{x} \in \mathscr{X}$ but not identification $\mathbf{H}$

$$
(i) \xrightarrow{(\mathrm{I})}(i i) \xrightarrow{(\mathrm{R})}(i i i)
$$

## Degrees of Freedom

$\triangleright \mathscr{H}=O P W(\Lambda) \subset \mathscr{L}\left(\mathbb{C}^{L}\right)$ is a linear subspace of dimension $|\Lambda|$.
$\triangleright$ Assume $\mathscr{X} \subset \mathbb{C}^{L}$ is a linear subspace of dimension $K$.
$\triangleright$ Counting degrees of freedom, we must have

$$
\begin{equation*}
|\Lambda|+K \leq L \quad \text { and } \quad|\Lambda| \leq L \tag{1}
\end{equation*}
$$

as a necessary condition for exact recovery of $\mathbf{H} \in \mathscr{H}$ and $\mathbf{x} \in \mathscr{X}$.
$\triangleright$ Without identifying $\mathbf{H}$, we may get $K>L-|\Lambda| \Longrightarrow$ how?
$\triangleright$ When we get equality in (1)?

## Example: 1-Dimensional Signal Space

$\triangleright$ For a given (and known) $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$, we consider the operator Paley-Wiener space

$$
\mathscr{H}=O P W(\Lambda)=\left\{\mathbf{H}=\sum_{k, \ell \in \Lambda} \eta(k, \ell) \mathbf{M}^{\ell} \mathbf{T}^{k}: \eta(k, \ell) \in \mathbb{C}\right\} \cdot \subset \mathscr{L}\left(\mathbb{C}^{L}\right)
$$

$\triangleright$ Assume a 1-dimensional signal space $\mathscr{X}=\operatorname{span}\{\mathbf{v}\} \subset \mathbb{C}^{L}$ for some $\mathbf{v} \in \mathbb{C}^{L}$.
$\triangleright$ For $x=u \mathbf{v} \in \mathscr{X}$, with $u \in \mathbb{C}$, the received signal is

$$
\mathbf{y}=\mathbf{H}(\mathbf{x}+\mathbf{c})=u \mathbf{H}(\mathbf{v})+\mathbf{H}(\mathbf{c})=\left.\mathbf{G}(\mathbf{c})\right|_{\wedge} \boldsymbol{\eta}+\left.u \mathbf{G}(\mathbf{v})\right|_{\wedge} \boldsymbol{\eta}
$$

$\triangleright$ Separation of $\mathbf{H x}$ and $\mathbf{H c}$ from $\mathbf{y} \Longrightarrow \operatorname{span}\{\mathscr{H} \mathbf{c}\} \cap \operatorname{span}\{\mathscr{H} \mathbf{v}\}=\{0\} \Longrightarrow \operatorname{rang}\left[\left.\mathbf{G}(\mathbf{c})\right|_{\wedge}\right] \perp \operatorname{rang}\left[\left.\mathbf{G}(\mathbf{v})\right|_{\wedge}\right]$.

For rang $\left[\left.\mathbf{G}(\mathbf{v})\right|_{\wedge}\right]=\operatorname{span}\{\mathbf{a}\}$ is a 1-dimensional subspace, we have the following result.

## Theorem

Let $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ with $1 \leq|\Lambda| \leq L-1$ and let $\mathbf{a} \in \mathbb{C}^{L} \backslash\{0\}$ be arbitrary.
There exists a vector $\mathbf{c} \in \mathbb{C}^{L}$ such that the $L \times|\Lambda|+1$ matrix $\left[\left.\mathbf{G}(\mathbf{c})\right|_{\Lambda}, \mathbf{a}\right]$ has full rank.

## General Construction

$\triangleright$ The set of all time-frequency shifts $\left\{\mathbf{M}^{\ell} \mathbf{T}^{k}\right\}_{\ell, k=0}^{L-1}$ can be separated into $L+1$ commutative subgroups

$$
\begin{aligned}
\mathscr{G}_{s} & =\left\{\mathbf{M}^{2 r k} \mathbf{T}^{k}: k=0,1, \ldots, L-1\right\}, \quad s=0,1, \ldots L-1 \\
\mathscr{G}_{L} & =\left\{\mathbf{M}^{k}: k=0,1, \ldots, L-1\right\}
\end{aligned}
$$

$\triangleright$ Each commutative subgroup $\mathscr{G}_{s}$ posses a set of common eigenvectors (i.e. the chirp sequences)

$$
\mathbf{e}_{s}(\ell): \ell=0,1, \ldots, L-1
$$

which forms on orthogonal basis for $\mathbb{C}^{L}$.

## Theorem

Let $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ such that there exists an $s \in\{0,1, \ldots, L\}$ so that $\Lambda \subset \mathscr{G}_{s}$.
There exists a subspace $\mathscr{X}=\operatorname{span}\left\{\mathbf{e}_{s}(1), \ldots, \mathbf{e}_{s}(K)\right\} \subset \mathbb{C}^{L}$ of dimension $K=L-|\Lambda|$ and a $\mathbf{c} \in \mathbb{C}^{L}$ such that

$$
\operatorname{span}\{\mathscr{H} \mathbf{c}\} \cap \operatorname{span}\{\mathscr{H} \mathscr{X}\}=\{0\}
$$

and such that Conditions (I) and (R) are satisfied.

## Maximum Size of Data Subspace

$\triangleright$ Given a subspace $\mathscr{X}$ of dimension $K$, it is desirable that the dimension of $\operatorname{span}\{\mathscr{H} \mathscr{X}\}$ is again $K$.
$\triangleright$ For for a one-dimensional $\mathscr{X}=\operatorname{span}\{\mathbf{v}\}$, we have $\operatorname{dim}(\operatorname{span}\{\mathscr{H} \mathscr{X}\})=\operatorname{rank}\left(\left.\mathbf{G}(\mathbf{v})\right|_{\wedge}\right)$.

Question: Given a support set $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ of $\mathscr{H}=\operatorname{OPW}(\Lambda)$ with $|\Lambda| \leq L$.
What is the minimum rank of $\left.\mathbf{G}(\mathbf{v})\right|_{\Lambda}$ for $\mathbf{v}$ varying in $\mathbb{C}^{L}$ ?

## Theorem

Let $L \geq 3$ be an odd integer and $\Lambda \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L}$ with $|\Lambda| \leq L$. Then

$$
\min _{\mathbf{v} \in \mathbb{C}^{\leq} \backslash\{0\}} \operatorname{rank}\left(\left.\mathbf{G}(\mathbf{v})\right|_{\Lambda}\right) \leq N(\Lambda)
$$

with

$$
N(\Lambda)=1+\min _{s \in\{0,1, \ldots, L\}} \min \left\{|I|: I \subset \mathbb{Z}_{L} \times \mathbb{Z}_{L} \text { with } \Lambda \subset I+\mathscr{G}_{s}\right\},
$$

and where $I+\mathscr{G}_{s}=\left\{x+y: x \in I, y \in \mathscr{G}_{s}\right\}$.

## Example with Unknown Channel Support

## Channel Model

$\triangleright$ Time continuous Time-Varying-Linear SISO Channel $\mathrm{H}: L^{2}(\mathbb{R}) \rightarrow L^{2}(\mathbb{R})$ :

$$
g(t)=(\mathrm{H} f)(t)=\iint_{\mathbb{R} \times \mathbb{R}} \eta_{\mathrm{H}}(\tau, v) \cdot f(t-\tau) \mathrm{e}^{\mathrm{i} 2 \pi v t} \mathrm{~d} v \mathrm{~d} \tau=\iint_{\mathbb{R} \times \mathbb{R}} \eta_{\mathrm{H}}(\tau, v)\left(\mathrm{M}_{v} \mathrm{~T}_{\tau} f\right)(t) \mathrm{d} v \mathrm{~d} \tau
$$

$\triangleright$ Time discrete Time-Varying-Linear SISO Channel H:L $L^{2}(\mathbb{R}) \rightarrow L^{2}(\mathbb{R})$ :

$$
g(t)=(\mathrm{H} f)(t)=\sum_{k=0}^{K-1} \sum_{\ell=0}^{M-1} \eta_{\mathrm{H}}(k, \ell) \cdot f(t-k \Delta \tau) \mathrm{e}^{\mathrm{i} 2 \pi \ell \Delta v t}=\sum_{k=0}^{K-1} \sum_{\ell=0}^{M-1} \eta_{\mathrm{H}}(k, \ell) \cdot f(t-k T) \mathrm{e}^{\mathrm{i} \frac{2 \pi}{T} \ell t}
$$

$\triangleright$ Rectification of the channel support region: with $\Delta \tau=T$ and $\Delta v=\frac{1}{T L}$ for some prime $L \geq 5$.


## Transmitted Signal

$\triangleright$ Transmitted signal: Delta train followed by a guard interval

$$
f(t)=\sum_{m=0}^{2 L-1} x_{m} \delta(t-m T) \quad \text { with } \quad x_{m}= \begin{cases}\text { data symbol } & : 0 \leq m \leq L-1 \\ 0 & : L \leq m \leq 2 L-1\end{cases}
$$

$\triangleright$ Received signal:

1. Sampling at rate $1 / T: \quad g_{n}=(\mathrm{H} f)(n T), n=0,1, \ldots, 2 L-1$.
2. Add two consecutive blocks of length $L: \quad y_{n}=g_{n}+g_{n+L}, n=0,1, \ldots, L-1$.
$\triangleright$ Write in vector form:

$$
\begin{equation*}
\mathbf{y}=\mathbf{H} \mathbf{x}=\sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta(k, l) \mathbf{M}^{\ell} \mathbf{T}^{k} \mathbf{x} \tag{2}
\end{equation*}
$$

with $\mathbf{x}=\left(x_{0}, \ldots, x_{L-1}\right)^{\mathrm{T}}$ and $\mathbf{y}=\left(x_{0}, \ldots, x_{L-1}\right)^{\mathrm{T}}$ and with $\mathbf{H} \in \operatorname{OPW}(\Lambda)$.

Remark: Using a periodic weighted delta train

$$
f(t)=\sum_{p \in \mathbb{Z}} \sum_{m=0}^{2 L-1} x_{m} \delta(t-m T-p L T)
$$

results in a similar expression as in (2) but requires periodic data.

## Transmission Scheme

$\triangleright$ Aim: We want to transmit a message $\boldsymbol{\gamma}=\left\{\gamma_{1}, \ldots, \gamma_{Q}\right\} \subset \mathbb{C}$ of size $Q \leq L$ over the channel $\mathbf{H} \in \operatorname{OPW}(\Lambda)$ and recover $\boldsymbol{\gamma}$ at the receiver without knowing $\mathbf{H}$ and without knowing the channel support $\wedge$ in advance.
$\triangleright$ Data symbols: The data symbols a sum of a pilot and the actual message

$$
\mathbf{x}=\mathbf{c}+\sum_{q=1}^{Q} \gamma_{q} \mathbf{e}_{q}
$$

with the pilot signal $\mathbf{c}$ which is chosen to be an Alltop sequence and $\mathbf{e}_{q}$ are particular chirp sequences

$$
\mathbf{c}(n)=\frac{1}{\sqrt{L}} \exp \left(\mathrm{i} \frac{2 \pi}{L} n^{3}\right) \quad \text { and } \quad \mathbf{e}_{m L+r}(n)=\frac{1}{\sqrt{L}} \exp \left(\mathrm{i} \frac{2 \pi}{L}\left[r+m n+r n^{2}\right]\right), \quad n \in \mathbb{Z}_{L}
$$

$\triangleright$ Received signal:

$$
\begin{equation*}
\mathbf{y}=\mathbf{H} \mathbf{x}=\sum_{\ell=0}^{L-1} \sum_{k=0}^{L-1} \eta(k, l) \mathbf{M}^{\ell} \mathbf{T}^{K} \mathbf{x}=\mathbf{G}(\mathbf{x}) \boldsymbol{\eta}=\mathbf{G}(\mathbf{c}) \boldsymbol{\eta}+\mathbf{U s}=\boldsymbol{\Phi}\binom{\boldsymbol{\eta}}{\mathbf{s}} \tag{3}
\end{equation*}
$$

with a $L \times 2 L^{2}$ measurement matrix $\boldsymbol{\Phi}=[\mathbf{G}(\mathbf{c}), \mathbf{U}]$ and a sparse vector $\mathbf{s}=f(\boldsymbol{\eta}, \boldsymbol{\gamma})$.
$\triangleright$ Recovery of the message: Solve CS problem (3), then determine $\boldsymbol{\gamma}=g(\boldsymbol{\eta}, \mathbf{s})$.

## Compressive Sampling Problem

$\triangleright$ Compressive Sampling Problem of size $L \times 2 L^{2}$

$$
\mathbf{y}=\boldsymbol{\Phi}\binom{\boldsymbol{\eta}}{\mathbf{s}} \quad \text { with } \quad\left|\operatorname{supp}(\boldsymbol{\eta}, \mathbf{s})^{\mathrm{T}}\right| \leq(1+Q)|\Lambda|
$$

Lemma
The coherence of the measurement matrix $\boldsymbol{\Phi}=[\mathbf{G}(\mathbf{c}), \mathbf{U}]$ is upper bounded by $\mu(\Phi) \leq \frac{2}{\sqrt{L}}$.
Remark: Welsh bound is $\frac{1}{\sqrt{L+1}}$.

## Theorem

Let $\mathbf{H} \in O P W(\Lambda) \subset \mathscr{L}\left(\mathbb{C}^{L}\right)$ be an unknown channel with unknown support set $\Lambda$ where $L \geq 5$ is a prime number. For

$$
Q \leq \frac{\sqrt{L}}{4|\Lambda|}-1
$$

any message $\boldsymbol{\gamma} \in \mathbb{C}^{Q}$ can be transmitted over $\mathbf{H}$ and recovered at the receiver.

## Numerical Experiment

- $L=307$
- OMP
- average over 100 channels (random support, Gaussian coefficients)
- red line: $1 \%$ rel. error
- simulation much better than bound



## Summary

## Current and Future Work

$\triangleright$ Identification of SISO and MIMO TVL Channels:

- Identification of stochastic channels and stochastic sequences
- Conditions on support of the covariance of the spreading function $\eta$.
- Linear side constraints in terms of covariance.
$\triangleright$ Transmission over Unidentified Channels:
- Stochastic encoding, RIP
- Scheme which maximum transmission rate
- Continuous time setting


## Related Publications

围 D．G．Lee，G．E．Pfander，V．Pohl，W．Zhou＂Identification of Channels with Single and Multiple Inputs and Outputs under Linear Constraints，＂Linear Algebra Appl．，vol． 581 （Nov．2019）， 435 － 470.

圁 D．G．Lee，G．E．Pfander，V．Pohl＂Signal transmission through an unidentified channel，＂ 13th Intern．Conf．on Sampling Theory and Applications（SampTA），Bordeaux，France，July 2019.

目 A．Kaplan，D．G．Lee，V．Pohl＂Message transmission through underspread time－varying linear channels，＂ 45th Intern．Conf．on Acoustics，Speech，and Signal Processing（ICASSP），Barcelona，Spain，May 2020

固 A．Kaplan，V．Pohl，D．G．Lee＂The Statistical Restricted Isometry Property for Gabor Systems＂，IEEE Statistical Signal Processing Workshop（SPP），Freiburg，Germany，June 2018， 45 － 49.

圁 D．G．Lee，G．E．Pfander，V．Pohl，W．Zhou＂Identification of multiple－input multiple－output channels under linear side constraints＇，＂43rd Intern．Conf．on Acoustics，Speech，and Signal Processing（ICASSP），Calgary，Canada， April 2018， 3889 － 3893.

